

# Recursive Data Types in HOL (temp extensions)

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## 1. INTRODUCTION

```
ML  
new_theory '051';;  
new_parent '042';;
```

```
ML  
load ('042',false);;
```

```
ML  
z_loading := false;;
```

```
ML  
z_loading := true;;
```

## 2. CONSTRUCTING EXISTENTIAL WITNESSES

For each type a witness must be constructed.

To do this the types must be processed in an appropriate order. A type is ready to have a witness constructed under the following conditions:

### *Disjoint Union*

A disjoint union is ready if any of its components has a witness.

### *Labelled Product*

A labelled product is ready if all of its components has a witness.

### *Void and Primitive components*

always have witnesses. The witnesses are constructed during the initialisation phase.

The code in this section does not merely construct witnesses. In the process of constructing the witnesses it constructs a plan for a proof that the witnesses satisfy the predicates.

The structure of this plan helps an understanding of the code and is documented by its ML type as follows.

```

type component      =   Prim_W    of term    |
                       Comp_W    of string;;

type trace          =   Prod_t    of component list |
                       Union_t   of int # component;;

typeabbrev witness  =   term # int # trace;;

typeabbrev witnesses =   (string # witness) list;;

```

The following function when given a witness, the predicate it is a witness for, and an object of type *witnesses* will return the theorem that the witness is indeed a witness.

```

let   REX_TAC (witness,depth,tr) witnesses = THENL[
  EXISTS_TAC witness;
  top_beta;
  EXISTS_TAC (nterm int);
  REWRITE_TAC[PRIM_REC_THM];
  TRACE_TAC witnesses tr]
where TRACE_TAC witnesses = fun
  Prod_t cl .

```

```

where   top_beta  =   CONV_TAC BETA_CONV
and     depth_beta =   CONV_TAC (DEPTH_CONV BETA_CONV);;

```

### 3. NON-RECURSIVE EXISTENCE PROOFS

#### 3.1. Primitive Types

#### 3.2. Unions

```

ML
let void_witproof rep_type =
  let [dom; cod] = snd (dest_type rep_type)
  in  "λv: ^rep_type.v=(list_mk_pfun [[: ^dom,μx: ^cod.T])";;

let witproof ((_, rep_type, terms,_,_),_,_,results,_) = fun
  P t  .fst (snd (snd (assoc t terms))),TRUTH      |
  T n  .snd (assoc n results),TRUTH                |
  V    .void_witproof rep_type,TRUTH;;

let witproofs g = map (witproof g);;

```

```

ML
let u_witness ((_,u,_,l),_,_,_) t =
  ("(inject ^ (dom_val l 1) ^t): ^u" );;

```

```

ML
let uex_tac wit1 thm wit2 = EVERY [
  EXISTS_TAC wit1;
  CONV_TAC BETA_CONV;
  DISJ1_TAC;
  EXISTS_TAC wit2;
  REWRITE_TAC [];
  ACCEPT_TAC thm];;

let uex_thm pred l thm wit2 =
  let goal = ([], "∃x. ^pred x")
  and wit1 = "inject ^ (dom_val l 1) ^wit2"
  in  TAC_PROOF(goal, uex_tac wit1 thm wit2);;

let do_nonrec_U_ex g (gn,s) =
  let ((_,rep_type,_,n),_,_,d,_) = g
  in  let pred          = snd (assoc gn d)
  and witness,proof = witproof g (hd s)
  in  (gn, witness, uex_thm pred n proof witness);;

```

### 3.3. Products

```

ML
let mk_product_ex g gn t = "t",TRUTH;;

```

```

ML
let do_nonrec_P_ex g (gn,s) =
  (gn,
   mk_product_ex g gn (map (witproof g) s));;

```

### 3.4. Non-recursive types

```

ML
let do_nonrectype_exists ad (name, st) =
  case st of   Union u   .do_nonrec_U_ex ad (name, map snd u)   |
             Prod p    .do_nonrec_P_ex ad (name, map snd p);;

```

```

ML
let exist_proved u v = true;;

let nonrectypes_estep (g,u,d,r,k) =
  let ok, not_ok = partition ((exist_proved u) o snd) u
  in   ok = [] => failwith 'nonrectypes_estep - finished' |
      (g,
       d @ ok,
       not_ok,
       r,
       k @ (map (do_nonrectype_exists (g,u,d,r,k)) ok));;

```

```

ML
let nonrectype_exists (g,u,d,r) =
  careful_repeat_f ['nonrectypes_estep - finished'] nonrectypes_estep (g,u,d,r,[]);;

```

## 4. CONTROL

```

ML
let new_rectypes = make_predicates o do_initiation;;

```

## **5. THE THEORY 051**